

Melocoton

A Program Logic for Verified Interoperability Between OCaml and C

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Michael Sammler, Lars Birkedal, Derek Dreyer

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MAX PLANCK INSTITUTE
FOR SOFTWARE SYSTEMS



AARHUS
UNIVERSITY
DEPARTMENT OF COMPUTER SCIENCE

Multi-Language Programs Are Everywhere



Python

C

Fortran



C++

Rust

JavaScript

OpenSSL
Cryptography and SSL/TLS Toolkit

C

Bindings for:

- Rust
- Python
- OCaml
- Go
- ...

The Goal: Verifying Multi-Language Programs

How do we

verify functional correctness

of programs written in

different languages?



Single-Language Functional Correctness

Hoare Logic for simple imperative languages.
Separation Logic for modularity and aliasing.

Multi-Language Functional Correctness





Multi-Language Functional Correctness

Existing work on Semantics and Logical Relations.

How do we prove functional correctness of individual, potentially unsafe programs?

A Multi-Language Program in OCaml and C

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C business logic

```
void hash_ptr(int * x) {  
    // Implemented in OpenSSL  
    // tedious to port to OCaml  
}
```


A Multi-Language Program in OCaml and C

OCaml business logic

```
let main () =  
  let r = ref 42 in  
  hash_ref r; (*written in C*)  
  print_int !r
```

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C glue code

```
value caml_hash_ref(value r) {  
  int x = Int_val(Field(r, 0));  
  hash_ptr(&x);  
  Store_field(r, 0, Val_int(x));  
  return Val_unit;  
}
```

A Multi-Language Program in OCaml and C

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OCaml glue code

```
external hash_ref  
  : int ref -> unit  
  = "caml_hash_ref"
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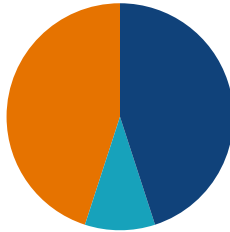
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A Schematic Multi-Language Program

Most multi-language programs look like this:

OCaml business logic
oblivious of C



C business logic
oblivious of OCaml

glue code

where the languages actually interact

We Need to Reason Language-Locally!

Our Contribution: Melocoton

λ_{ML+C} **Program Logic**

Glue Code Verification

λ_{ML+C} **Semantics**

Glue Code Semantics

Common Approach: program logic on top of semantics, **but**

- **Language Interaction:** new semantics and logic for glue code

Our Contribution: Melocoton



Common Approach: program logic on top of semantics, **but**

- **Language Interaction:** new semantics and logic for glue code
- **Language Locality:** embed existing semantics and logics

*simplified/idealized versions of OCaml and C

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Language Interaction: Different Views of the Same Data

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How is **OCaml data** accessed from **C glue code**?

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How is **OCaml data** accessed from **C glue code**?

High-level **OCaml values** are accessed..
..through a **low-level block representation**.

Language Interaction: Semantics

High-level **OCaml** value \sim_{ML} Low-level **block** representation

Language Interaction: Semantics

High-level **OCaml** value \sim_{ML} Low-level **block** representation

integers \sim_{ML} integers

booleans \sim_{ML} integers (0 or 1)

true \sim_{ML} *1*

Language Interaction: Semantics

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λ_{ML+C} Semantics

$\sigma : Heap_{ML}$

$\zeta : BlockHeap$

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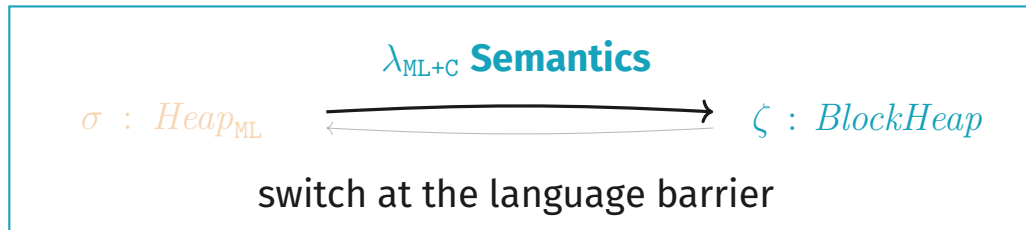
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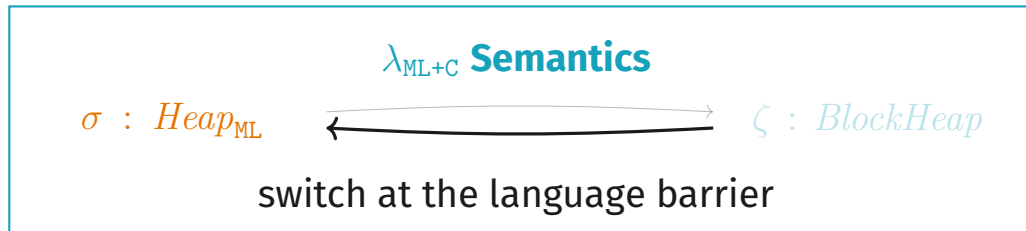
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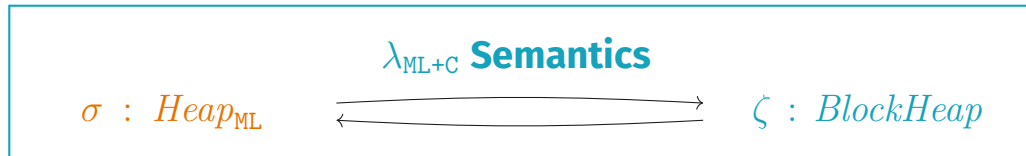
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Language Interaction: Program Logic, Take 1

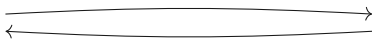


Language Interaction: Program Logic, Take 1

$\lambda_{\text{ML+C}}$ Program Logic

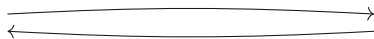
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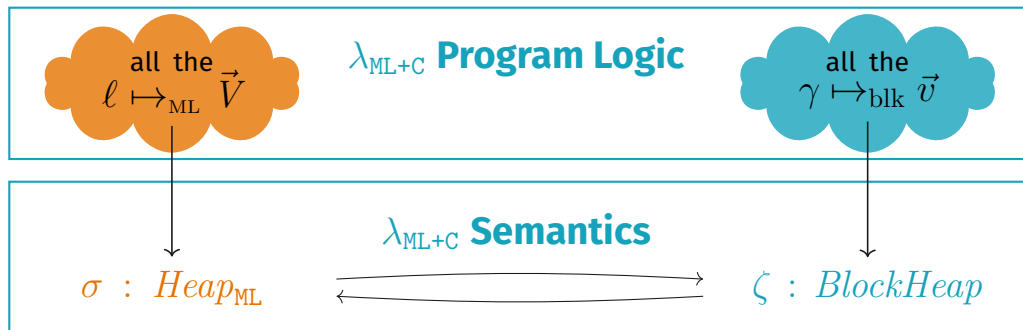


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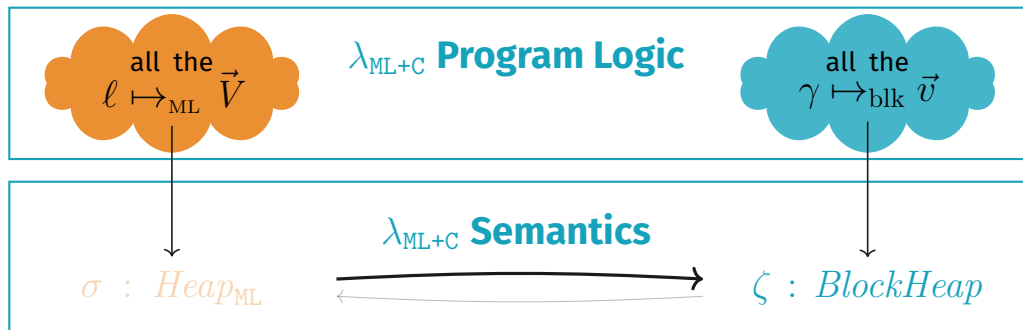
Language Interaction: Program Logic, Take 1

 $\ell \mapsto_{\text{ML}} \vec{V}$ $\lambda_{\text{ML+C}}$ **Program Logic** $\gamma \mapsto_{\text{blk}} \vec{v}$ $\sigma : \text{Heap}_{\text{ML}}$ $\lambda_{\text{ML+C}}$ **Semantics** $\zeta : \text{BlockHeap}$ 

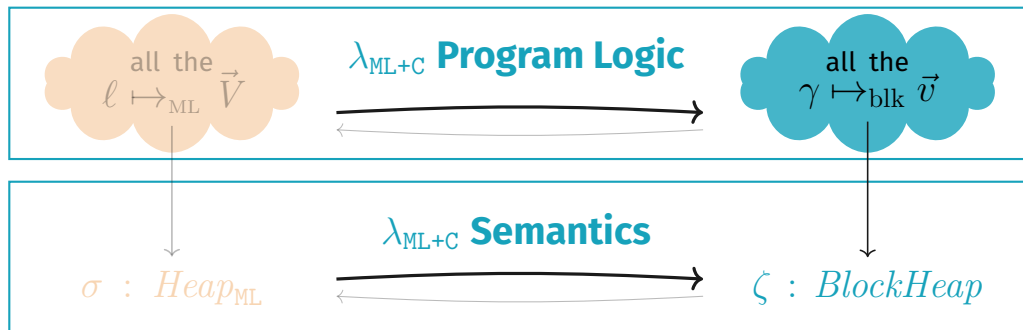
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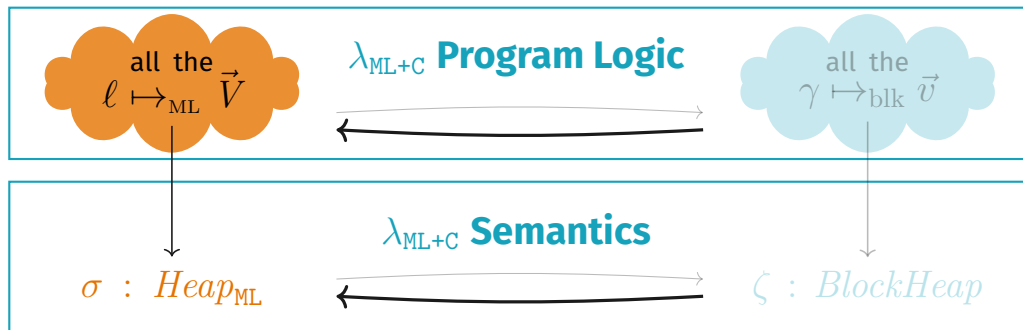
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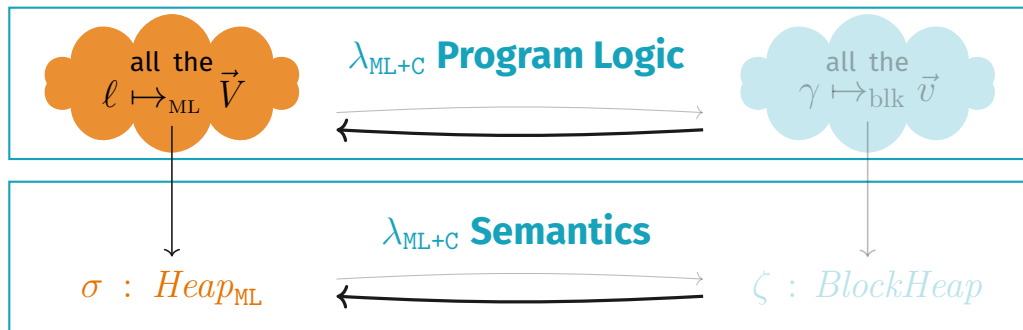
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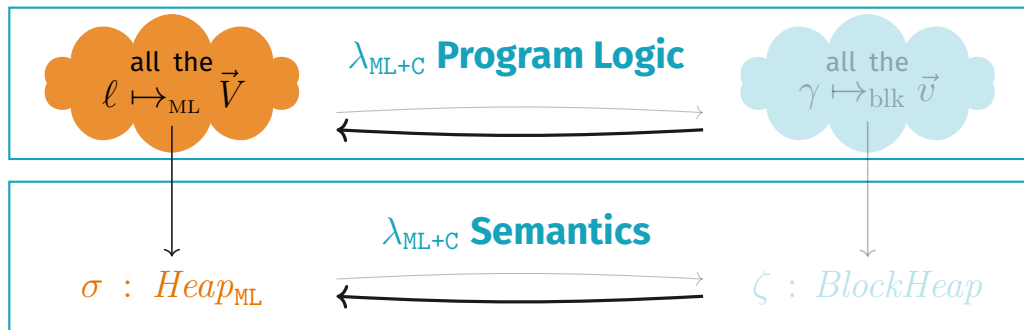


EXTCALL

{ all } C function body { all }

{ all } call into C { all }

Language Interaction: Program Logic, Take 1



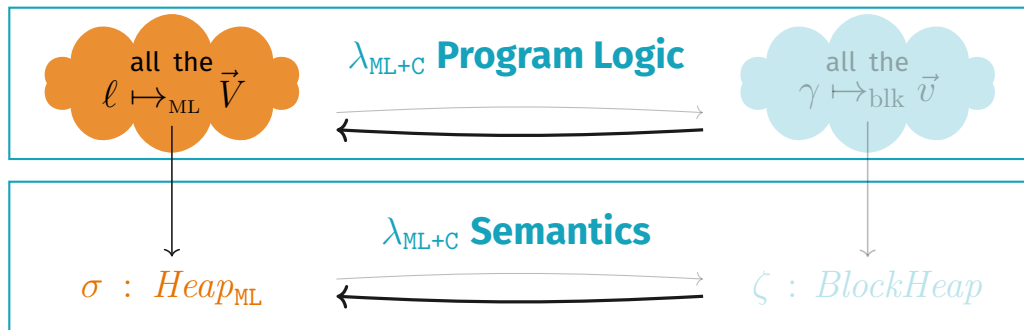
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$$\frac{\{ \text{all} \} \mathbf{C} \text{ function body } \{ \text{all} \}}{\{ \text{all} \} \text{ call into } \mathbf{C} \{ \text{all} \}}$$

FRAME

$$\frac{\{P\} e \{Q\}}{\{R * P\} e \{Q * R\}}$$

Language Interaction: Program Logic, Take 1



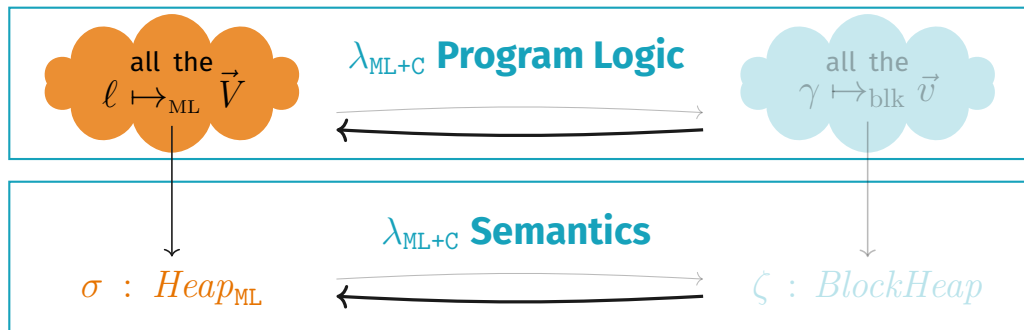
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Language Interaction: Program Logic, Take 1



EXTCALL

$\{\text{all}\} \mathbf{C}$ function body $\{\text{all}\}$

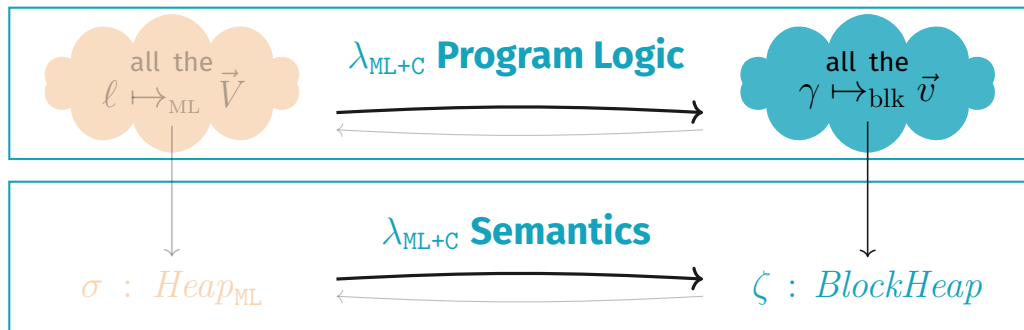
$\{\text{all}\}$ call into \mathbf{C} $\{\text{all}\}$

FRAME

$\{P\}$ call into \mathbf{C} $\{Q\}$

$\{\ell \mapsto_{ML} \vec{V} * P\}$ call into \mathbf{C} $\{Q * \ell \mapsto_{ML} \vec{V}\}$

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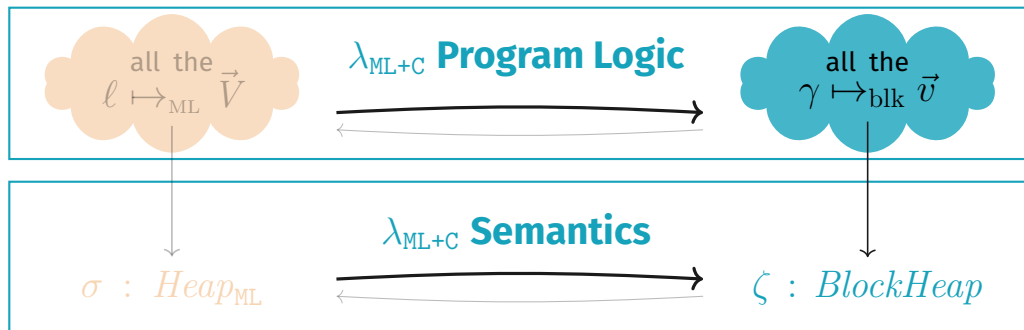
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Language Interaction: More Gradual Rules

OCaml *points-tos* remain valid when switching to **C**!

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$$l \mapsto_{\text{ML}} \vec{V} \quad l_1 \mapsto_{\text{ML}} \vec{V}_1$$

Language Interaction: More Gradual Rules

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$$l \mapsto_{\text{ML}} \vec{V} \quad \gamma_1 \mapsto_{\text{blk}} \vec{v}_1$$

Language Interaction: More Gradual Rules

OCaml points-tos *remain valid* when switching to **C!**

$$\gamma_2 \mapsto_{\text{blk}} \vec{v}_2$$

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View Reconciliation Rules for Converting On-Demand:

$$\begin{aligned} l \mapsto_{\text{ML}} \vec{V} &\equiv * \exists \gamma \vec{v}. \gamma \mapsto_{\text{blk}} \vec{v} * l \sim_{\text{ML}} \gamma * \vec{V} \sim_{\text{ML}} \vec{v} \\ \vec{V} \sim_{\text{ML}} \vec{v} * \gamma \mapsto_{\text{blk}} \vec{v} &\equiv * \exists l. l \mapsto_{\text{ML}} \vec{V} * l \sim_{\text{ML}} \gamma \end{aligned}$$

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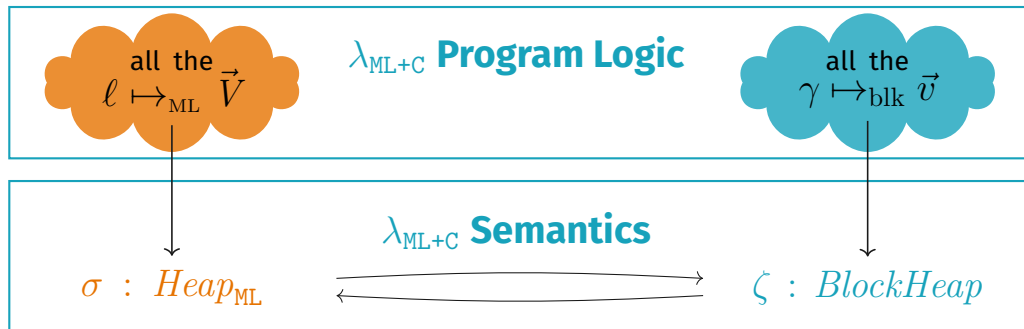
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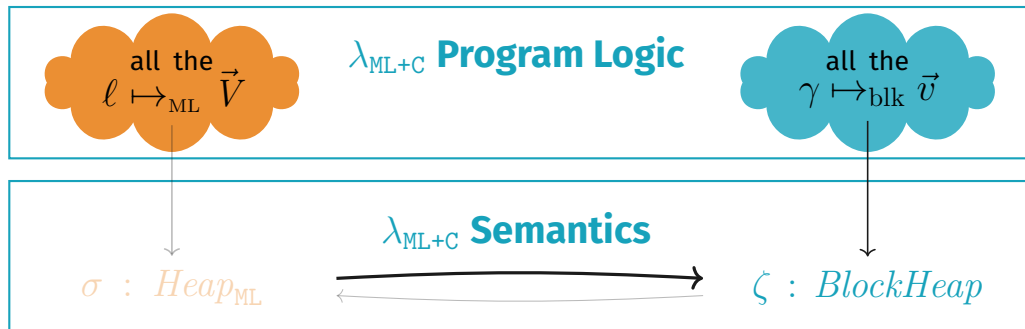
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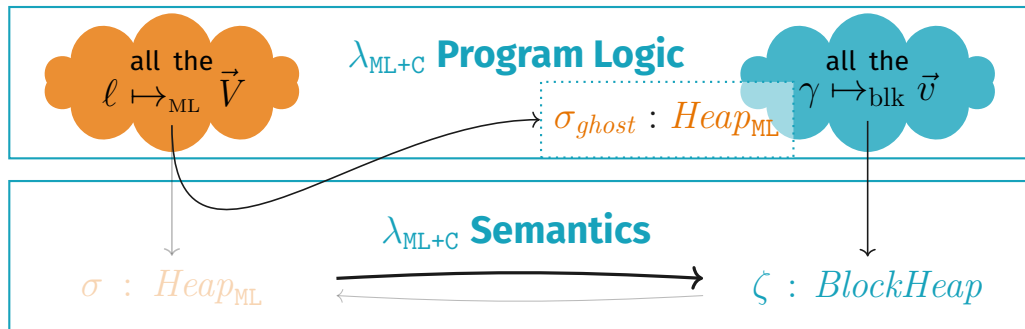
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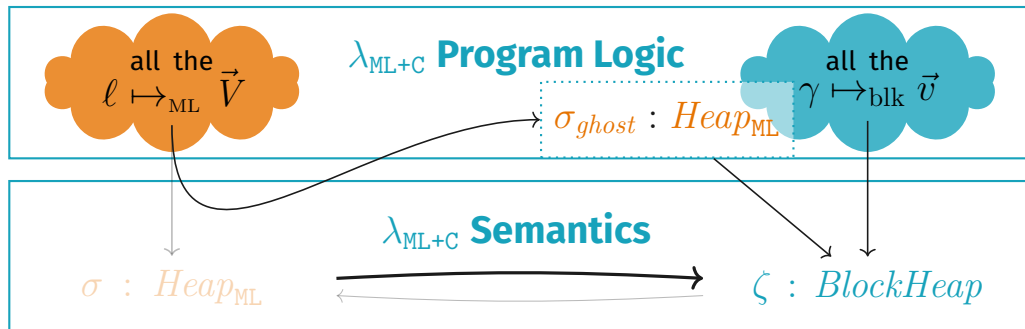
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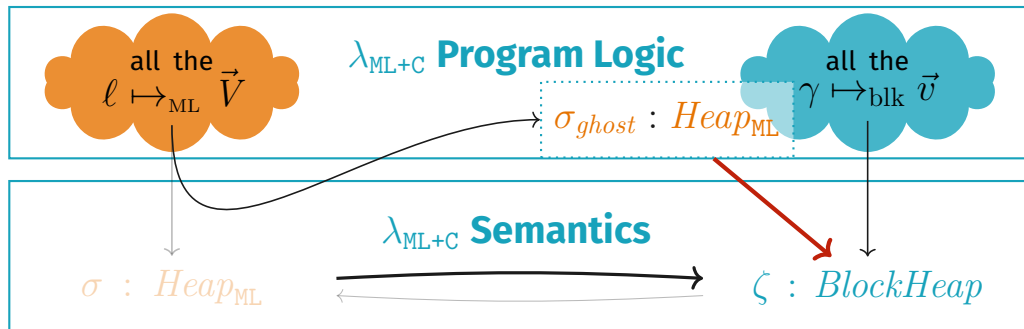
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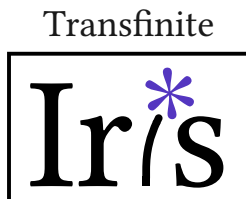
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More in the paper ...

- Language-local reasoning for **external calls**.
- Additional **OCaml FFI features**: garbage collection, registering roots, custom blocks, callbacks, etc.
- **Case studies** utilising all of these features.
- **Step-indexed logical relation** to prove OCaml type safety of external C functions.



Our Contribution: Melocoton

Language Locality: Embed Existing Languages



Language Interaction: View Reconciliation Rules

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<https://melocoton-project.github.io>